

# MTH314: Discrete Mathematics for Engineers

## Lecture 1: Propositional Logic

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11 January 2017

# Logistics

Section 2:

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Evaluation:

iClickers	15%+5% bonus	starting week 2
Midterm I	20%	on February 15
Midterm II	25%	on March 22
Final Exam	40%	final exam period

# Logistics

Online resources:

- d2l
- Discussion board on piazza.com

Your best resources to succeed in the class:

- 1
- 2
- 3

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# Let's get started!

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We have a statement  $P$  and assign it a value 1 if  $P$  is true and 0 if  $P$  is false.

Example:  $P$  : "I am 21 years old."

Truth table (depicting all the possibilities):

$$\begin{array}{c} P \\ \hline 0 \\ 1 \end{array}$$



## A more elaborate truth table.

$P$  : "I am 21 years old."

$Q$  : "I play soccer."

$P$	$Q$
0	0
0	1
1	0
1	1

Which row of the table corresponds to you?

# Negation

Statement: a declarative sentence that is either true or false. We will denote statements with capital letter, for example  $P$ ,  $Q$ ,  $R \dots$

The *negation*  $\neg P$  of a statement  $P$  is a statement that evaluates to true if  $P$  evaluates to false, and evaluates to false if  $P$  evaluates to true. (Some textbooks write  $\sim P$  or “NOT  $P$ ” instead.)

Example:

$P$  : “Earth is flat.”

$\neg P$  : “Earth is not flat.”

$P$	$\neg P$
0	1
1	0

# Conjunction and Disjunction

Conjunction (AND) is true if and only if both statements are true.

$$P \wedge Q$$

Disjunction (OR) of two statements is true if at least one of the statements is true.

$$P \vee Q$$

$P$  : "I am 21 years old."

$Q$  : "I play soccer."

$P \wedge Q$  : "I am 21 years old AND I play soccer."

$P \vee Q$  : "I am 21 years old OR I play soccer."

Remember: the Disjunction symbol points Down. And the symbol for Conjunction looks a bit like the "A" in AND.

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$P \wedge Q$  : "I am 21 years old AND I play soccer."

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Draw a truth table for these.

$P$	$Q$	$P \wedge Q$	$P \vee Q$
0	0		
0	1		
1	0		
1	1		

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# Implication and Equivalence

$P$  : "I am 21 years old."

$Q$  : "I play soccer."

$P \rightarrow Q$  : "IF I am 21 years old THEN I play soccer."

(IMPLICATION)

$P \leftrightarrow Q$  : "I am 21 years old IF AND ONLY IF I play soccer."

(DOUBLE IMPLICATION/EQUIVALENCE)

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## Unary and binary operators.

```
def BinaryAddition(x,y):  
    return x+y  
    print "I require two inputs."
```

```
def UnaryAddition(x):  
    return x+7  
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$$P \vee Q \vee R$$

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The conjunction, disjunction, implication and equivalence operator are binary, they need exactly two inputs.

$$(P \vee Q) \vee R$$

## Combining logical operators.

We evaluate  $\neg$  before the binary operators, that is:

$$\neg P \vee Q \equiv (\neg P) \vee Q,$$

as opposed to:

$$\neg(P \vee Q).$$

$P$	$\neg P$	$Q$	$\neg P \vee Q$	$\neg P \rightarrow Q$	$(\neg P \vee Q) \leftrightarrow (\neg P \rightarrow Q)$
1	0	1			
1	0	0			
0	1	1			
0	1	0			

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1	0	1	1	1	1
1	0	0	0	1	0
0	1	1	1	1	1
0	1	0	1	0	0

## Remarks.

Just as  $\neg$  comes before  $\wedge$  and  $\vee$ ,  
 $\wedge$  and  $\vee$  come before  $\rightarrow$  and  $\leftrightarrow$ .

Two logical statements for which their truth tables agree are called  
**logically equivalent.**

A logical statement that evaluates always to true is called a  
**tautology.**

A logical statement that evaluates always to false is called a  
**contradiction.**

$P$  is a tautology if and only if  $\neg P$  is a contradiction.

$P$  and  $Q$  are logically equivalent if and only if  $P \leftrightarrow Q$  is a  
tautology.

# Examples of tautologies

- $P \vee \neg P$
- $(P \rightarrow Q) \vee P$
- De Morgan's Laws:
  - $\neg(P \vee Q) \leftrightarrow \neg P \wedge \neg Q$
  - $\neg(P \wedge Q) \leftrightarrow \neg P \vee \neg Q$

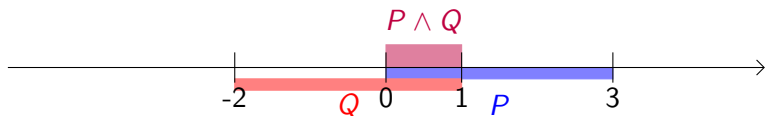
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  - $\neg(P \wedge Q) \leftrightarrow \neg P \vee \neg Q$

Example:

$P$  :  $x$  is a real number greater than 0 but smaller than 3.

$Q$  :  $x$  is a real number smaller than 1 but greater than -2.



## Examples of contradictions

- $P \wedge \neg P$
- $(P \rightarrow Q) \wedge (P \wedge \neg Q)$
- $P \vee Q \leftrightarrow \neg P \wedge \neg Q$
- $P \wedge Q \leftrightarrow \neg P \vee \neg Q$

## Some important math expressions

For some logical statements  $P$ ,  $Q$  :

“ $P$  is a **sufficient condition** for  $Q$ .”

“ $P$  is true **only if**  $Q$ .”

means  $P \rightarrow Q$  is a tautology.

“ $P$  is a **necessary condition** for  $Q$ .”

“ $P$  is true **if**  $Q$ .”

means  $Q \rightarrow P$  is a tautology.

“ $P$  is a **necessary and sufficient condition** for  $Q$ ”

“ $P$  is true **if and only if**  $Q$ ”

“ $P$  is equivalent to  $Q$ .”

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Example:

$P$  : “This is a Tesla Model X”

$Q$  : “This is a car.”

We have  $P \rightarrow Q$ , because any Tesla Model X is a car. However, not every car is a Tesla Model X, so while  $P$  is a sufficient condition for  $Q$ , it is not necessary.

“This is a Tesla Model X **only if** it is a car.”



# Can you come up with a sentence in English (rather than math) for each of these tautologies?

“ $\equiv$ ” reads as “is equivalent to”, “is the same as” or simply “is.”

- $\neg(\neg P) \equiv P$  (double negative)
- $P \leftrightarrow Q \equiv (P \rightarrow Q) \wedge (Q \rightarrow P)$
- $P \vee \text{false} \equiv P$
- $P \wedge \neg P \equiv \text{false}$
- $P \vee (Q \wedge R) \equiv (P \vee Q) \wedge (P \vee R)$  (distributive law)
- $P \vee Q \equiv Q \vee P$  (commutative law)
- $P \rightarrow (Q \rightarrow R) \equiv Q \rightarrow (P \rightarrow R)$

# What can we do with an implication?

Example:

$P$  : It is Monday.

$Q$  : There is no MTH314 class.

We know that  $P \rightarrow Q$  is true. What about the following statements?

**Converse:**  $Q \rightarrow P$

**Inverse:**  $\neg P \rightarrow \neg Q$

**Contrapositive:**  $\neg Q \rightarrow \neg P$

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Inverse:  ~~$\neg P \rightarrow \neg Q$~~

Contrapositive:  $\neg Q \rightarrow \neg P$  TRUE

# What is a proof?

Start with some statements you assume to be true and use the properties of logical operations to show that something else is true.

It's a bit like arguing a case in court. You start with some assumptions, like laws and evidence, and make a step-by-step, air-tight argument to convince the judge of your conclusion.

Example: Suppose you know that both  $P \rightarrow Q$  and  $P$  are true. Then  $Q$  is true. You can use a truth table to convince a picky judge.

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Then  $Q$  is true.

$P$	$Q$	$P \rightarrow Q$
1	1	1
1	0	0
0	1	1
0	0	1

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$P$  is true.

$P$	$Q$	$P \rightarrow Q$
1	1	1
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$P$  is true.

$P$	$Q$	$P \rightarrow Q$
1	1	1
1	0	1
0	1	0
0	0	1

# What is a proof?

An argument of this form:

$$\begin{array}{l} P \rightarrow Q, P \\ \therefore Q \end{array}$$

Is called “Modus Ponens.” (The  $\therefore$  symbol reads as “therefore.”)

We call  $P \rightarrow Q$  and  $P$  the *premises* and  $Q$  the *conclusion*.

You can similarly verify other arguments:

$$\begin{array}{l} P \rightarrow Q, \neg Q \\ \therefore \neg P \end{array}$$

What about

$$\begin{array}{l} P \rightarrow Q \\ \therefore Q \rightarrow P? \end{array}$$

# Proof by Contradiction

Proof by contradiction is of this form:

$\neg P \rightarrow \text{false}$

$\therefore P$

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$\neg P$  : There exists a Tesla Model X that is not a car.

If we can argue that  $\neg P$  is false, we conclude that  $P$  is true.

$P$	$\neg P$	false
1	0	0
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1	0	0
<del>0</del>	<del>1</del>	<del>0</del>

$\neg P \rightarrow \text{false}$